

EXERCISE #6

DATAFLOW REVIEW

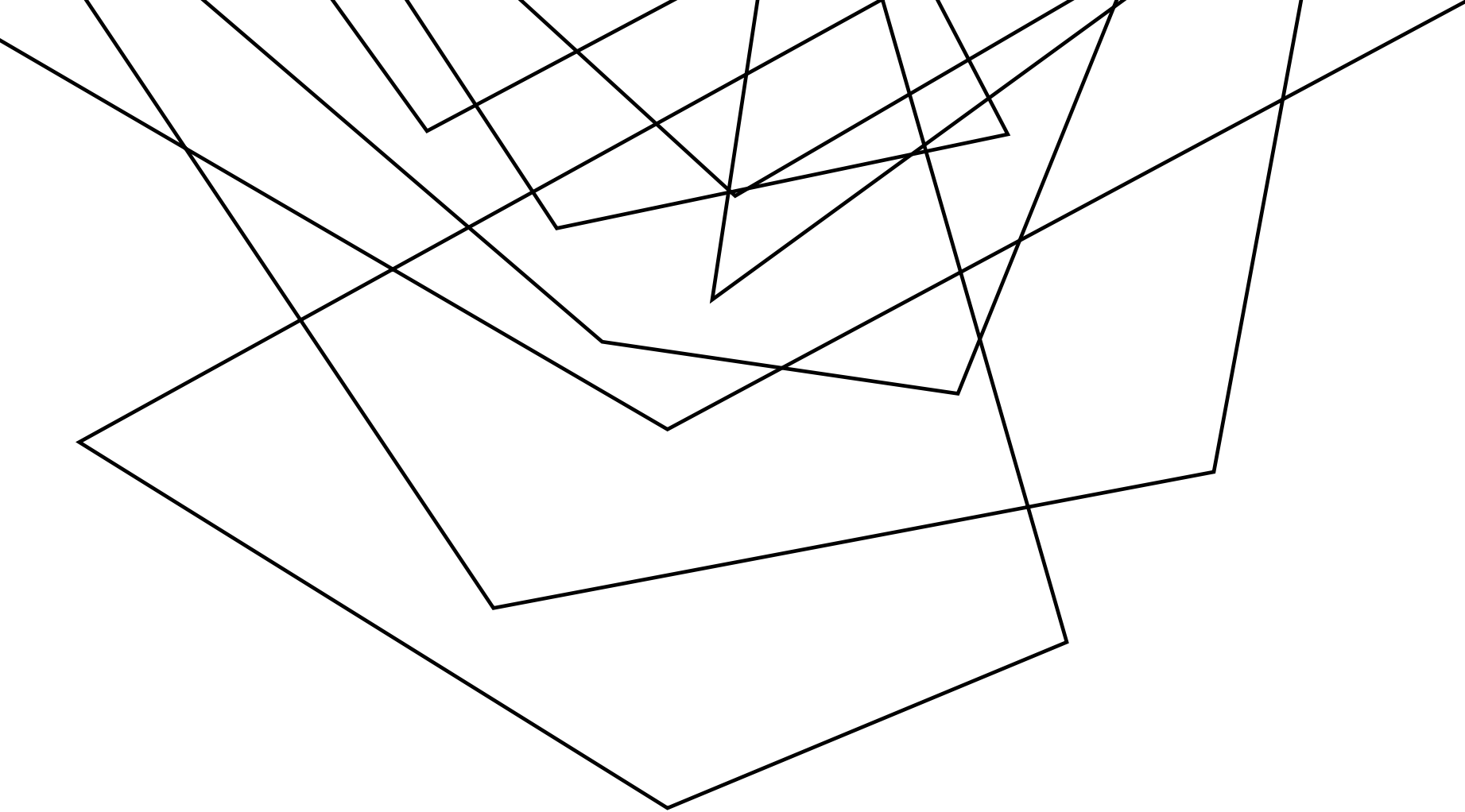
Write your name and answer the following on a piece of paper

- Show the control-flow graph of the following function and indicate the value-sets at the beginning and end of each basic block according to the method discussed in class

```
y = 0;
if ( g ) {
    x = 1;
    x += y;
} else {
    x = 3;
    if (g2) {
        x = y;
    }
    x = 4;
}
z = x;
```



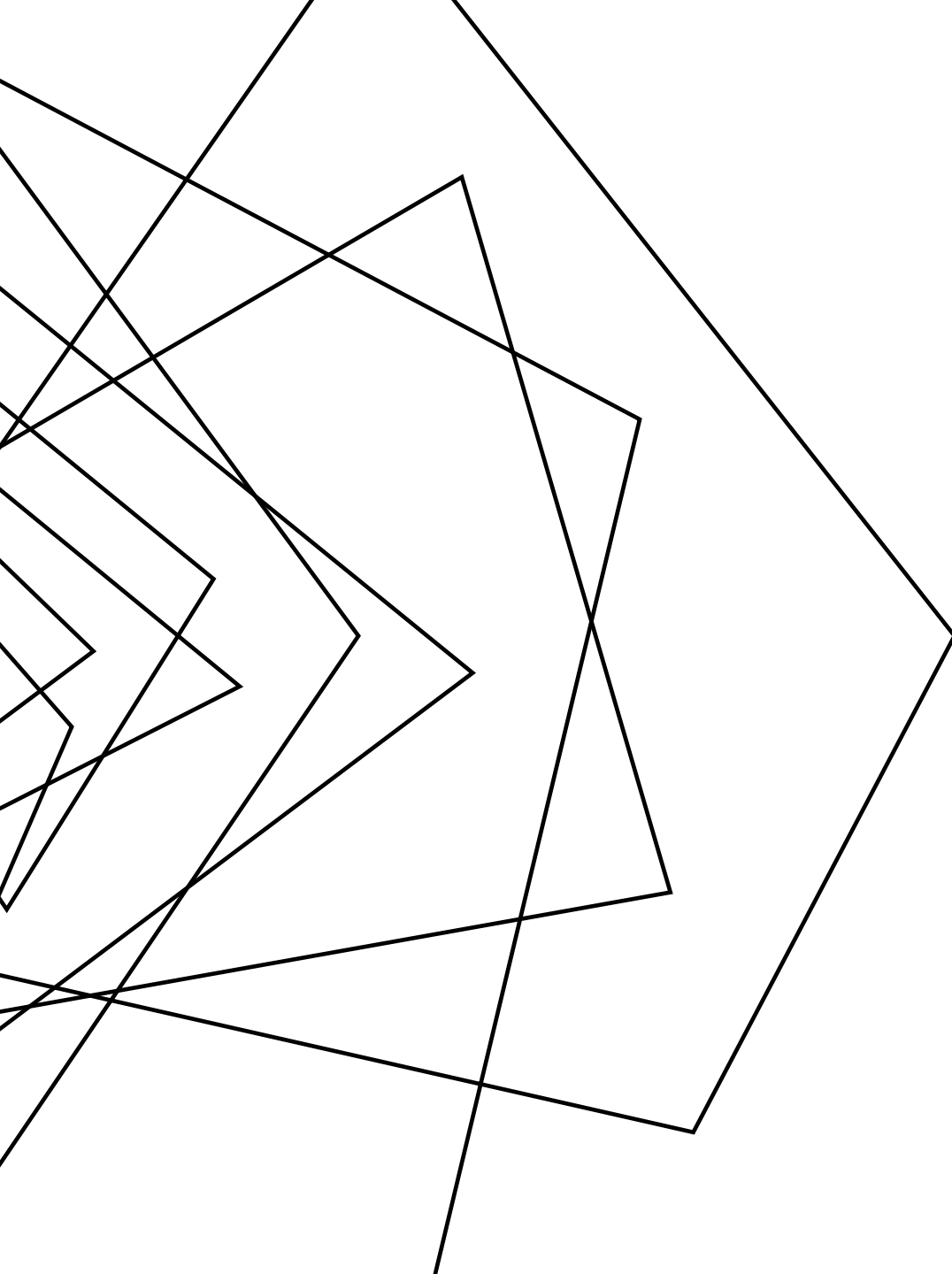
**ADMINISTRIVIA
AND
ANNOUNCEMENTS**



DATAFLOW FIXPOINTS

EECS 677: Software Security Evaluation

Drew Davidson



CLASS PROGRESS

EXPLORING A FORM OF STATIC ANALYSIS THAT SUMMARIZES HOW CONTROL AND DATA FLOWS ACROSS A PROGRAM

- MANIFEST A COMPLETE ANALYSIS BY DENOTING SETS OF ALL VALUES MEMORY MIGHT CONTAIN (**NB – THIS WILL END UP BEING CUMBERSOME!**)

LAST TIME: CONTROL-FLOW GRAPHS

REVIEW: STATIC ANALYSIS

SET RULES FOR PARTITIONING FUNCTIONS INTO BASIC BLOCKS

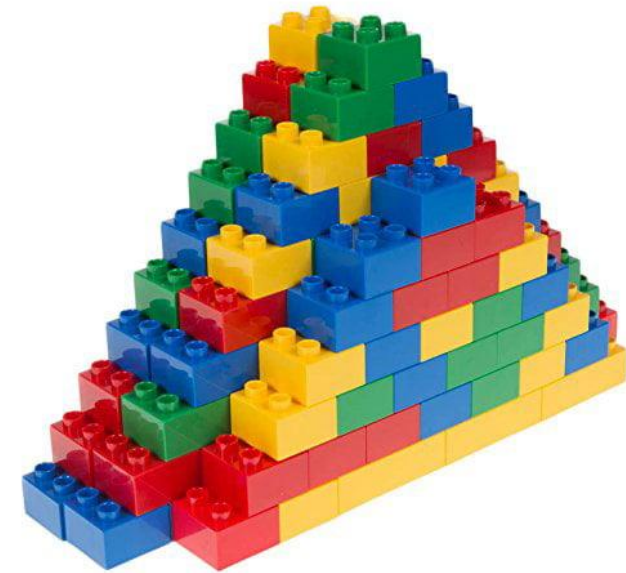
- Leader instructions
- Terminator instructions

CONNECTED THE BASIC BLOCKS INTO A CONTROL-FLOW GRAPH

- Edges indicate control flow transfers

EXPLORED THE “BIG IDEA” OF DATAFLOW ANALYSIS

- Treat each instruction as a transfer function
- Compose transfer functions to model blocks data flow
- Merge block effects to get function's data flow



Building something bigger out of basic blocks

LAST TIME: FLOW SENSITIVE VALUE SETS

REVIEW: DATAFLOW

ACCOUNT FOR PROGRAM FLOW, NOT PATHS

- Necessarily Over-approximating states

LET'S DO AN EXAMPLE

```
1. int x = 0;
2. int y = 0;
3. if (g == true) {
4.     x = 10;
5. }
6. if (g == false) {
7.     y = 1 / (x - 10);
8. }
9. return;
```

LOOPS ARE TOUGH TO HANDLE!

REVIEW: DATAFLOW ANALYSIS

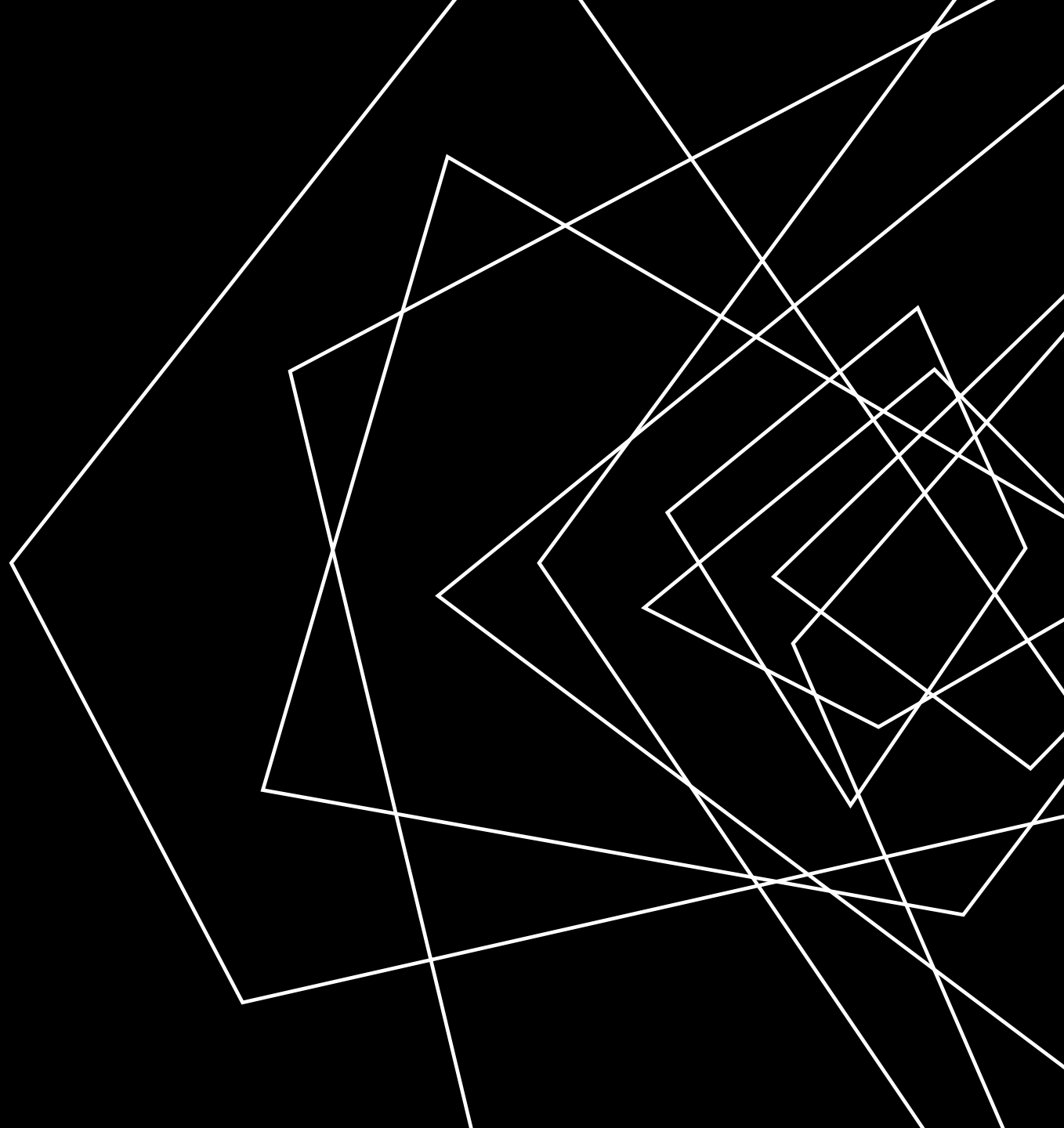
ISSUES WITH LOOPS

- Generate lots of paths
- Cyclic data dependency



LECTURE OUTLINE

- Breaking cyclic dependency
- Termination
- Handling large value sets



WHERE TO START ANALYSIS?

BREAKING CYCLIC DEPENDENCY

```

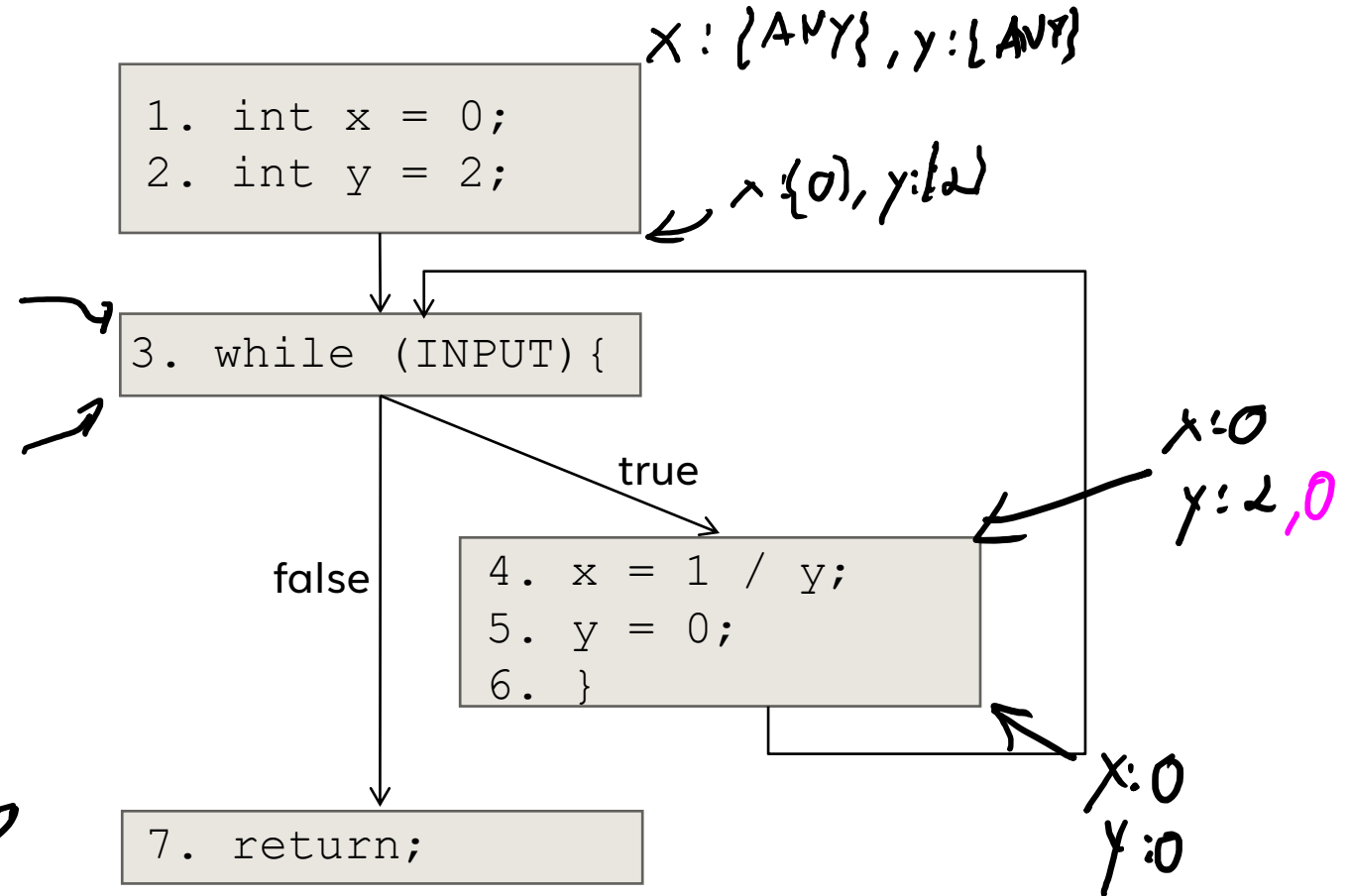
1. int x = 0;
2. int y = 2;
3. while (INPUT) {
4.   x = 1 / y;
5.   y = 0;
6. }
7. return;

```

x:0
y:2,0

x:0
y:2,0

x:0
y:2,0



CHAOTIC ITERATION

STATIC ANALYSIS: CONTROL FLOW GRAPHS

A WORKLIST ALGORITHM

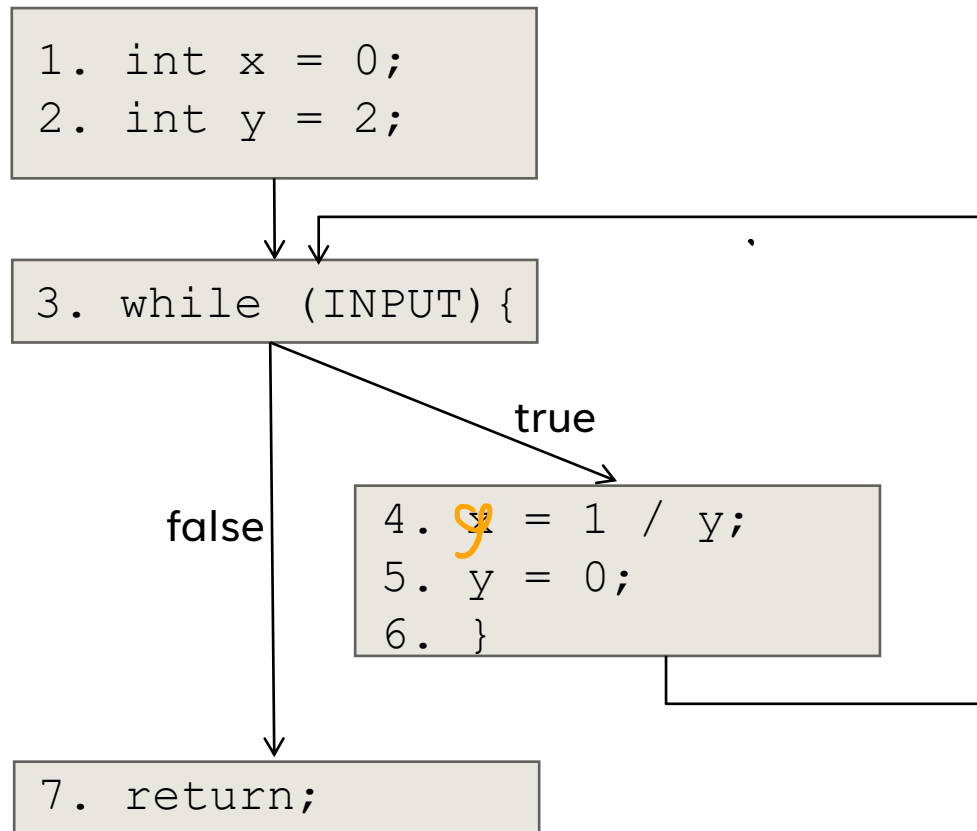
- Select the next worklist item in any order
- Necessarily assumes progress towards some goal



Surprisingly, not a band with merch at Hot Topic

WHERE TO START ANALYSIS?

STATIC ANALYSIS: CONTROL FLOW GRAPHS



$x: \{1\}, y: \{1\}$
 $B_{1-2} \quad x: \{1\} \quad y: \{1, 2\}$

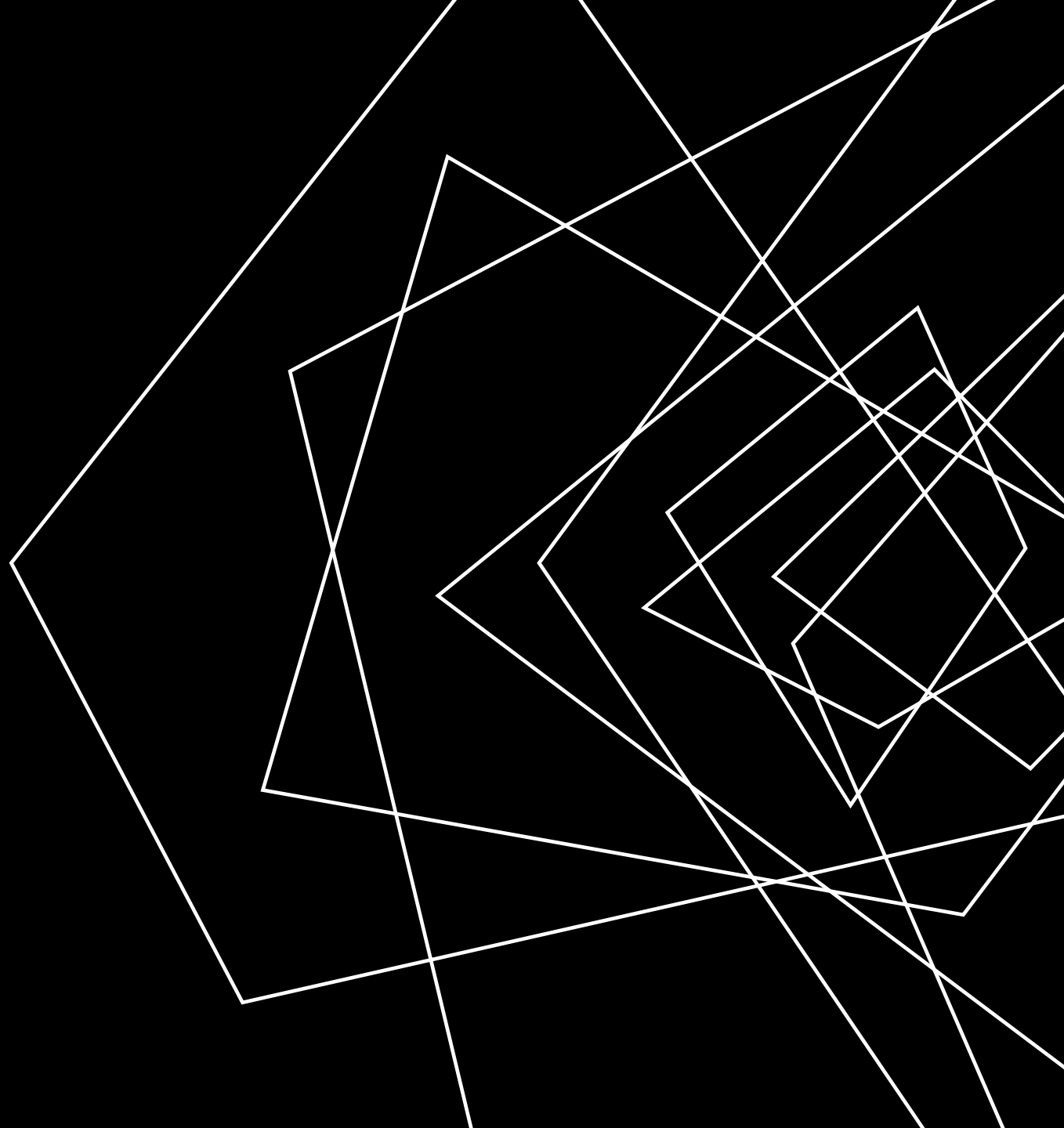
 $B_3 \quad x: \{1, 2\} \quad y: \{1, 2, 0\}$
 $x: \{1\} \quad y: \{1\}$

 $B_{4-6} \quad x: \{1, 2\} \quad y: \{1, 2, 0\}$

 $B_7 \quad x: \{1, 2\} \quad y: \{1\}$
 $x: \{1\} \quad y: \{1\}$

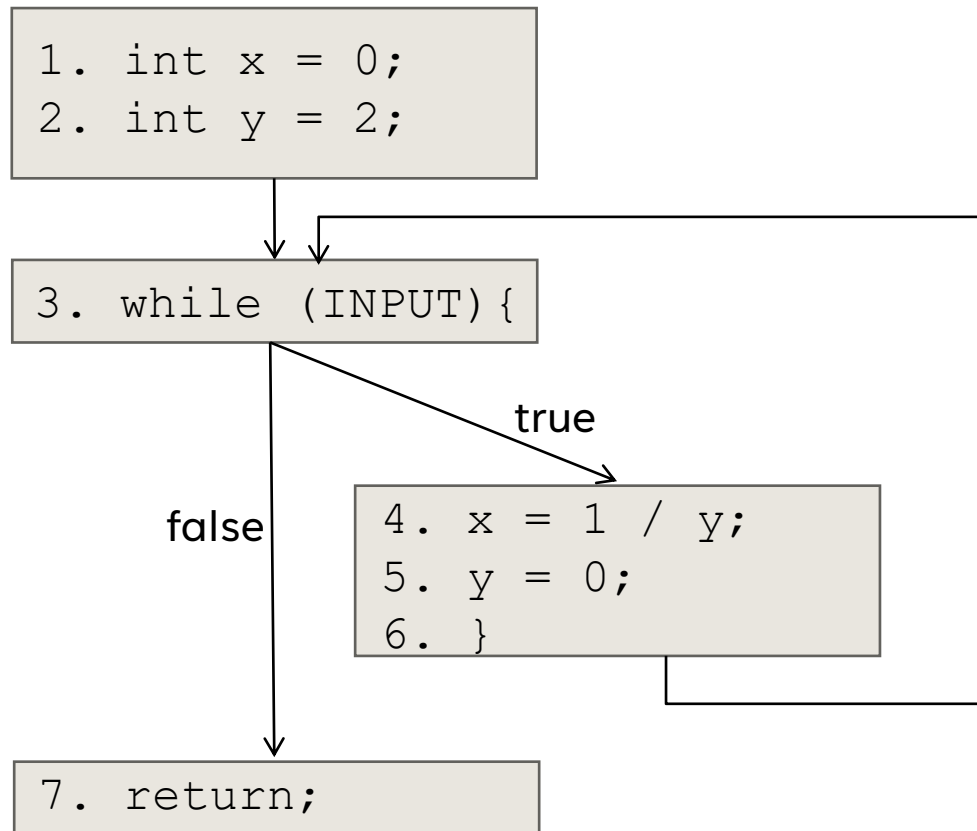
LECTURE OUTLINE

- Breaking cyclic dependency
- Termination
- Handling large value sets



WHERE TO STOP ANALYSIS?

STATIC ANALYSIS: CONTROL FLOW GRAPHS



ANALYSIS PROGRESS

STATIC ANALYSIS: CONTROL FLOW GRAPHS

ANALYSIS ENDS WHEN THE FACT SETS REACH *SATURATION*

- No additional elements will ever be added
- It sure would be nice if we could guarantee that this will happen!



When your fact sets couldn't possibly hold any more data

FIXED-POINTS

STATIC ANALYSIS: CONTROL FLOW GRAPHS

A FIXED-POINT (AKA FIXPOINT, FIXED POINT)

- A value that does not change under a given transformation

OUR VALUE-SET ANALYSIS WILL HAVE FACTS THAT REACH A FIXED-POINT

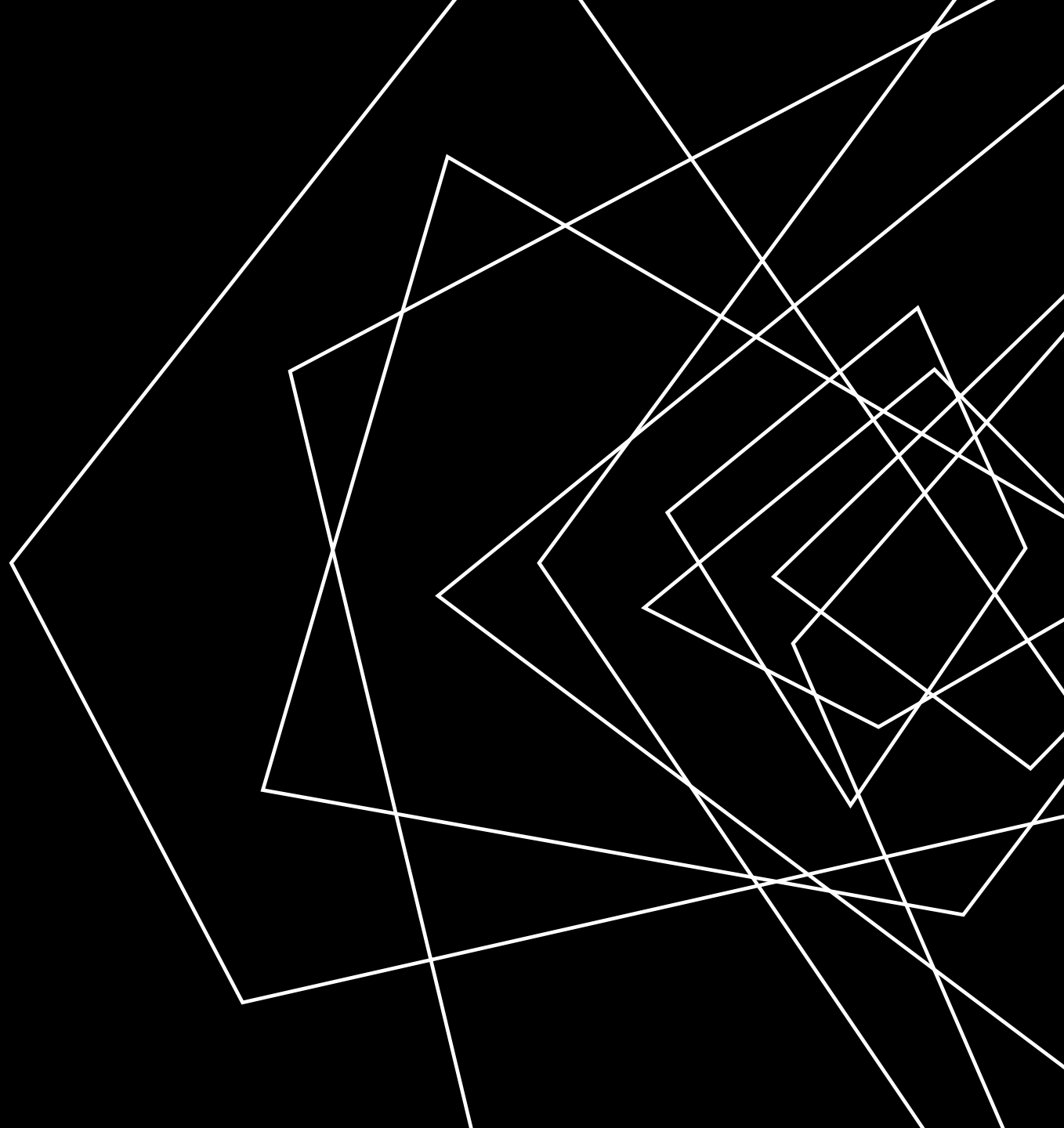
Why?

- Finite set of configurations over INT32s
- Data transforms only add data to fact sets

$$\begin{aligned} & \{ a, b, c \} \\ & \cup \\ & \{ a \} \\ & = \\ & \{ a, b, c \} \end{aligned}$$

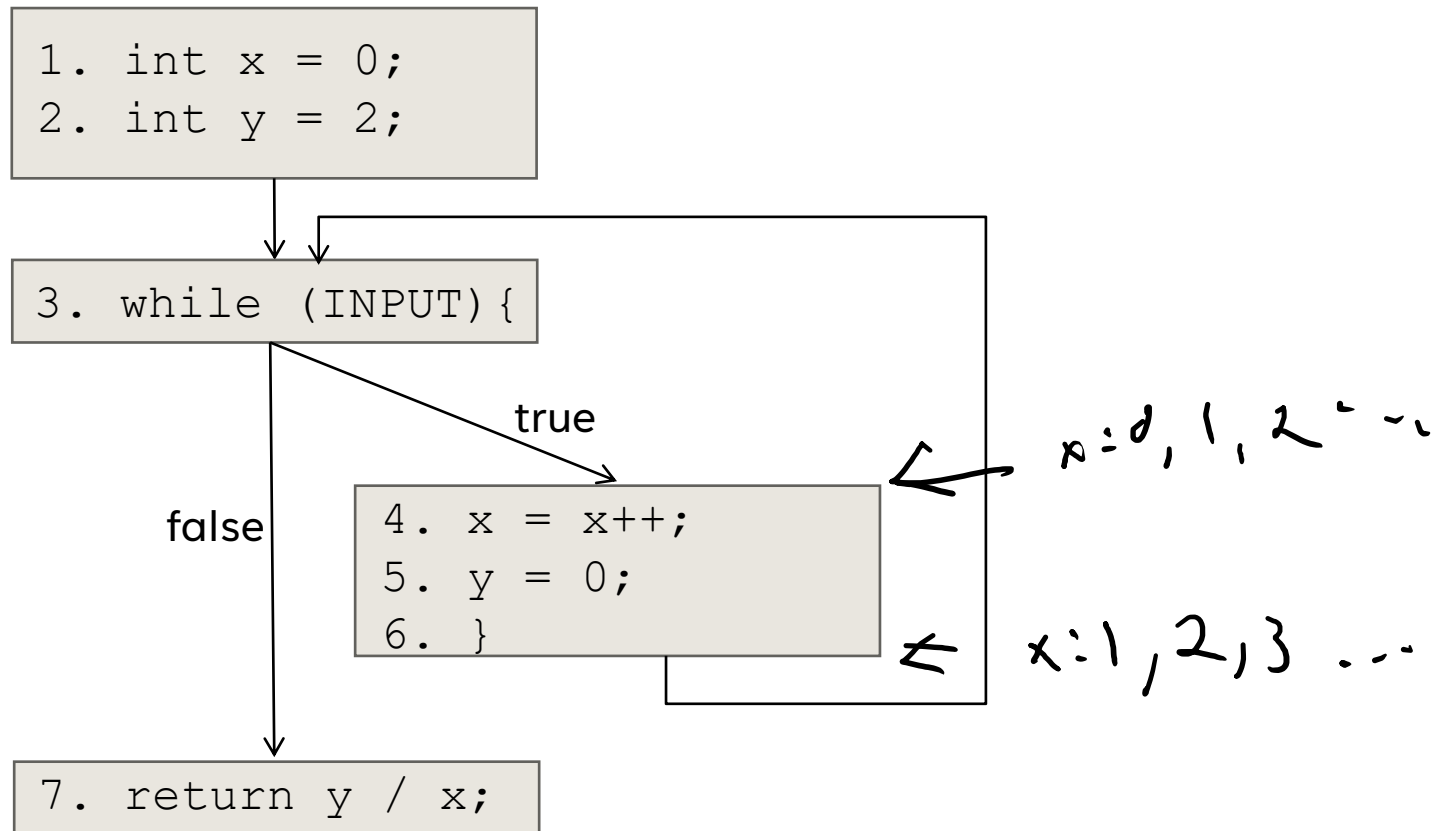
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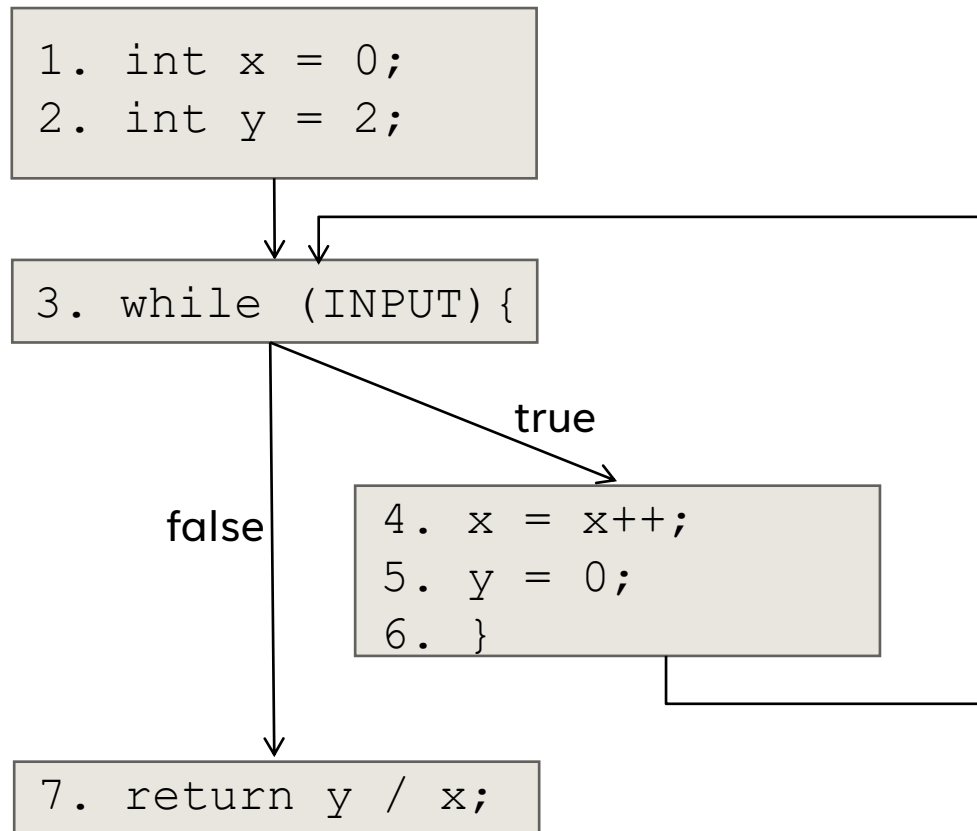
WHERE TO STOP THIS ANALYSIS?

ANALYSIS TERMINATION



WIDENING

ANALYSIS TERMINATION



ACCELERATE PROGRESS TOWARDS FIX-POINT

- Lots of (over-approximate) ways to do this
- 1 simple idea: if we hit N values, immediately change the fact set to “All integers”

$x: 0, 1, 2, 3 \Rightarrow x: ANY$

$x: 1, 2, \Rightarrow x: ANY$



LECTURE END!

*DESCRIBED SOME OF THE ISSUES AND FIXES
FOR DATAFLOW IN THE PRESENCE OF LOOPS*