

EXERCISE #14

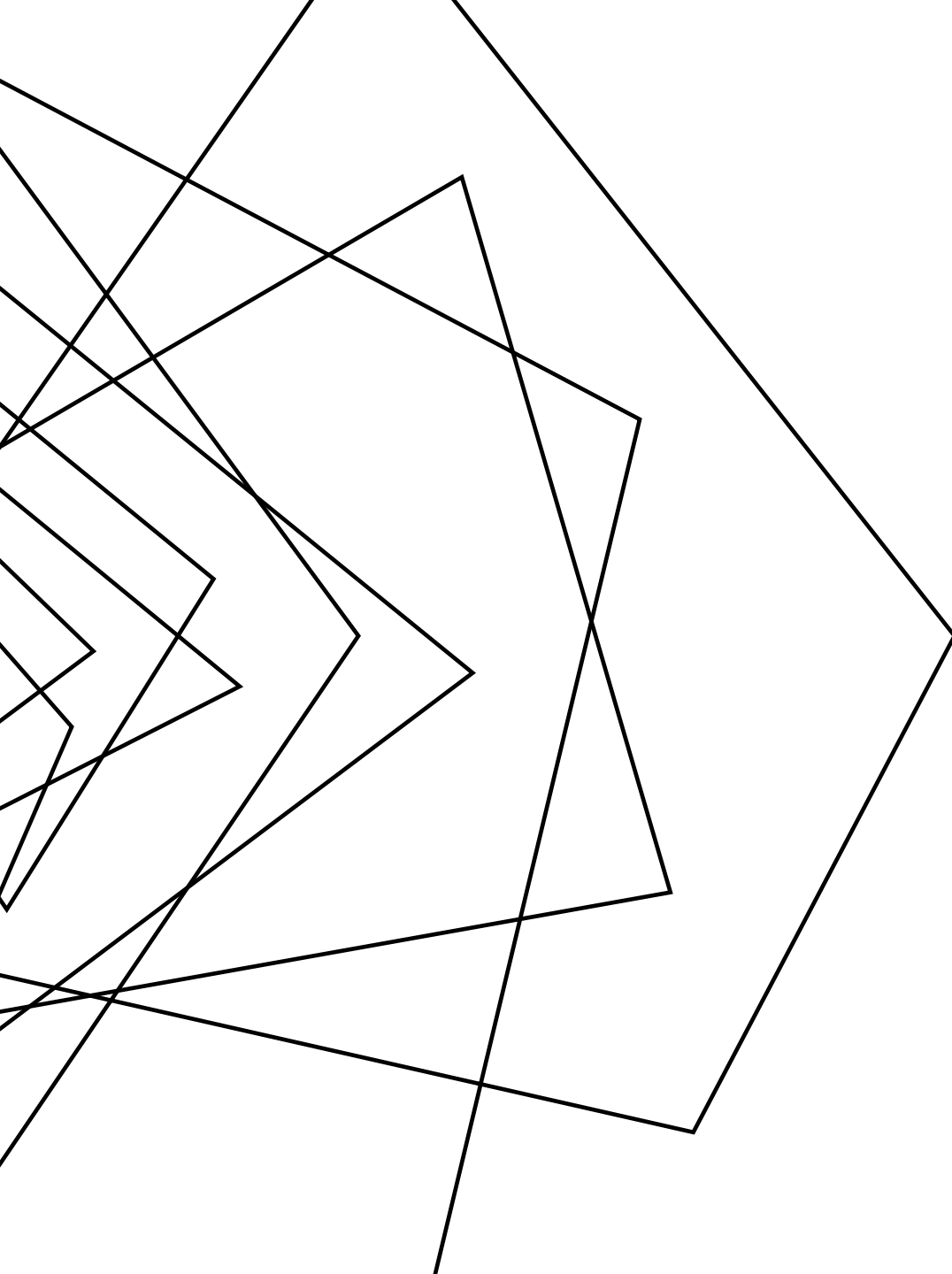
PRACTICAL INFOFLOW REVIEW

Write your name and answer the following on a piece of paper

Provide an instance of a program with an implicit information flow from a confidential source to a sink



**ADMINISTRIVIA
AND
ANNOUNCEMENTS**



CLASS PROGRESS

SHOWING SOME APPLICATIONS OF
STATIC DATAFLOW

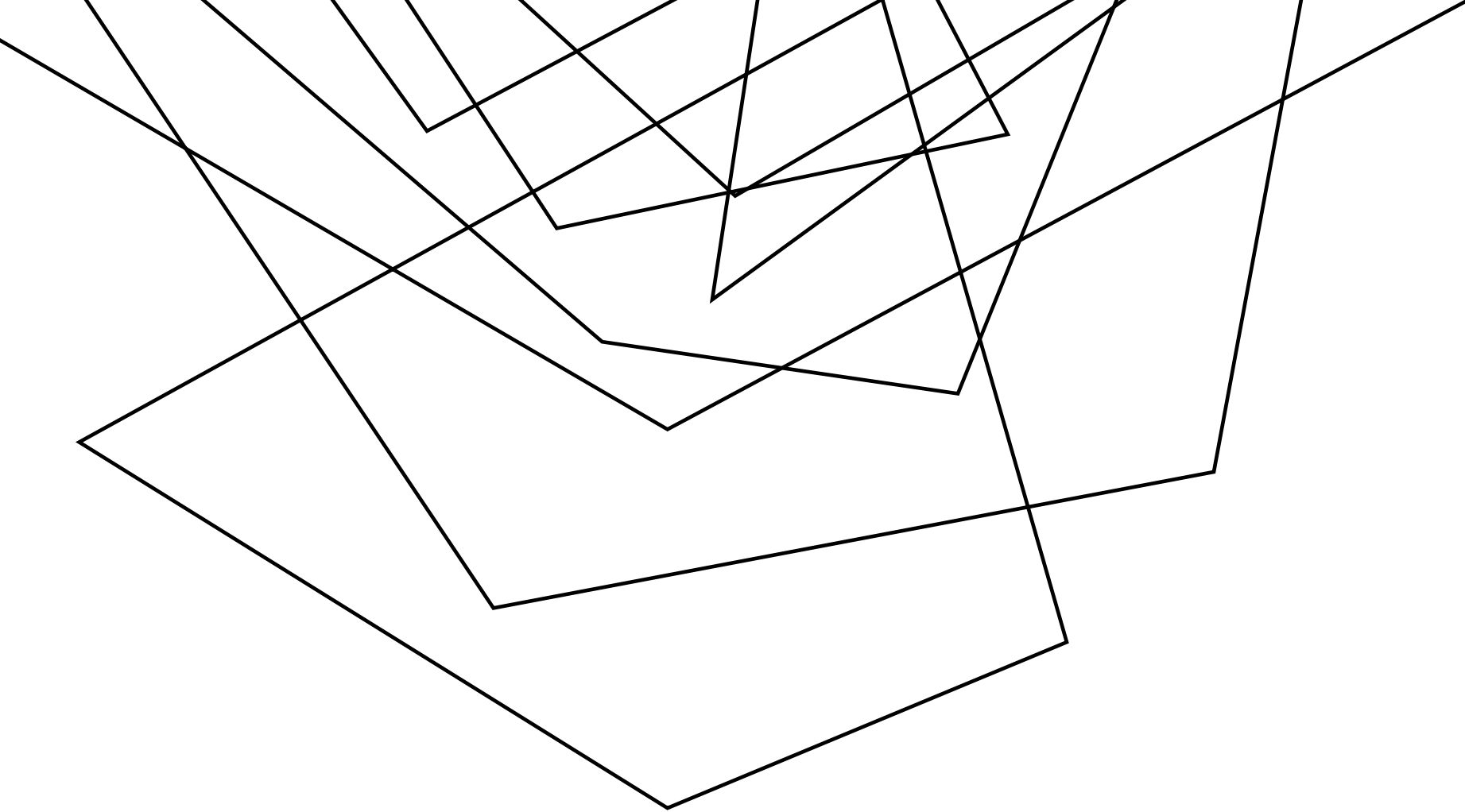
LAST TIME: DATAFLOW DEPLOYMENT

REVIEW: LAST LECTURE

USING DATAFLOW IN PRACTICAL CONTEXTS

- Ex. - Looking for initialized variables

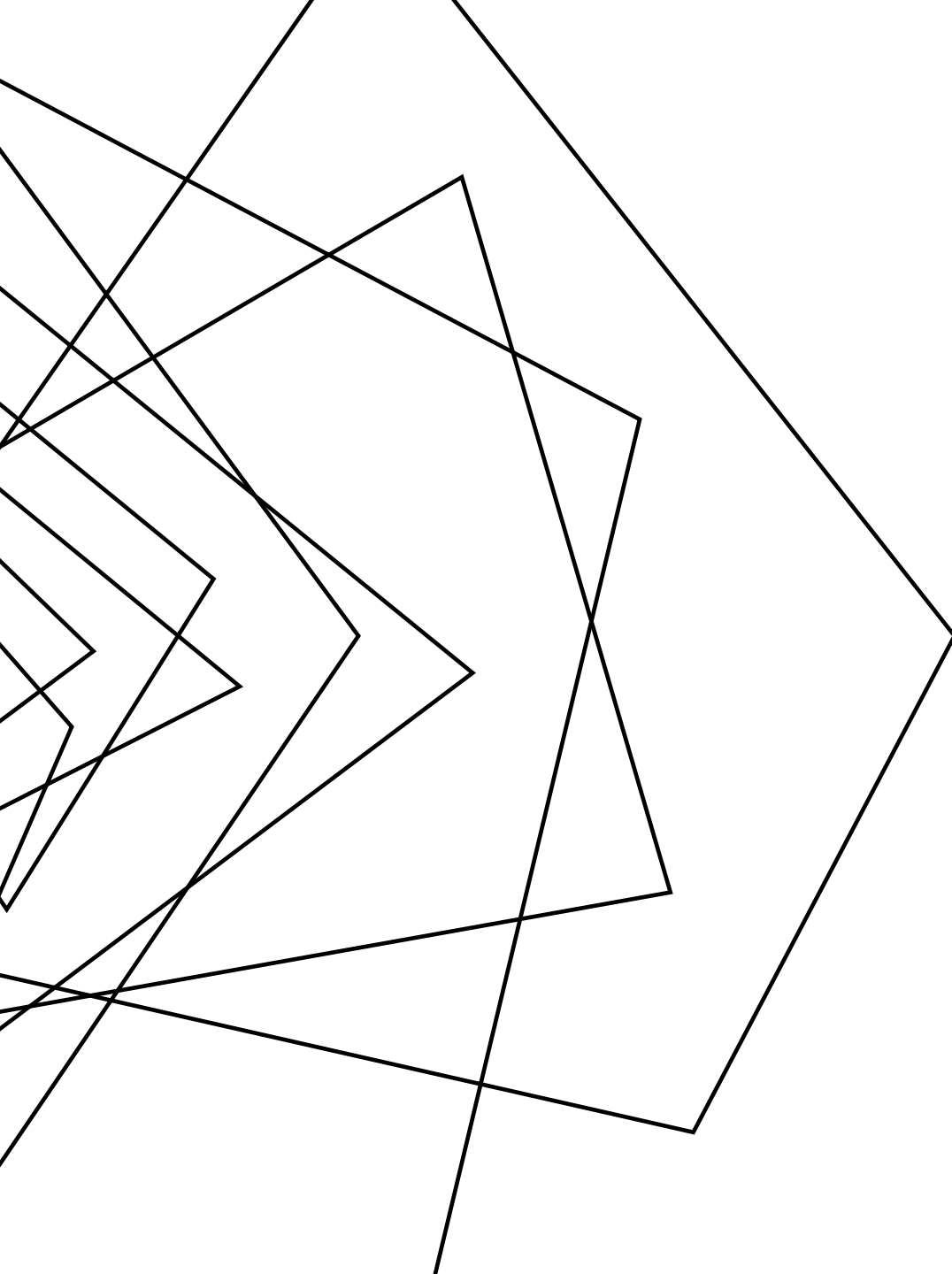




SIDE CHANNELS

EECS 677: Software Security Evaluation

Drew Davidson

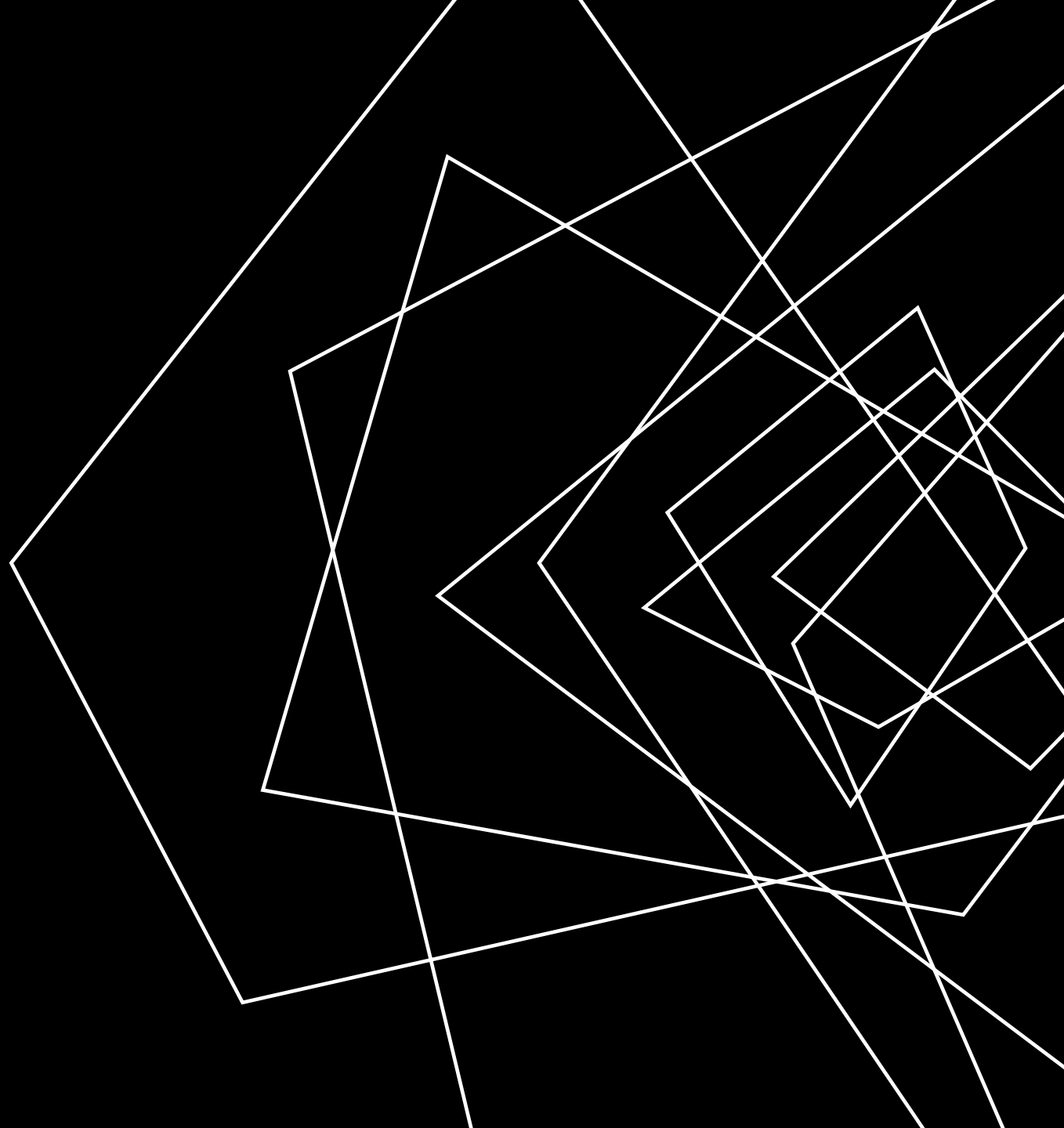


OVERVIEW

CONTEMPLATE OTHER WAYS THAT
SNEAKY DATA FLOWS CAN OCCUR

LECTURE OUTLINE

- Threat Models
- Side Channels - Overview
- Timing
- A dataflow approach



THINKING ABOUT ATTACKS

THREAT MODELS

THERE'S NO SUCH THING AS "ABSOLUTE SECURITY"

- It's always possible to come up with SOME (potentially wacky) scenario where the adversary can subvert a system

CONSIDER THE VARIOUS ATTACK CLASSES

- **Denial of Service:** Availability is compromised
- **Exfiltration:** Confidentiality policy is compromised
- **Compromise:** Integrity policy is compromised



A FRAMEWORK FOR ASSUMPTIONS

THREAT MODELS

A THREAT MODEL IS COMPOSED OF:

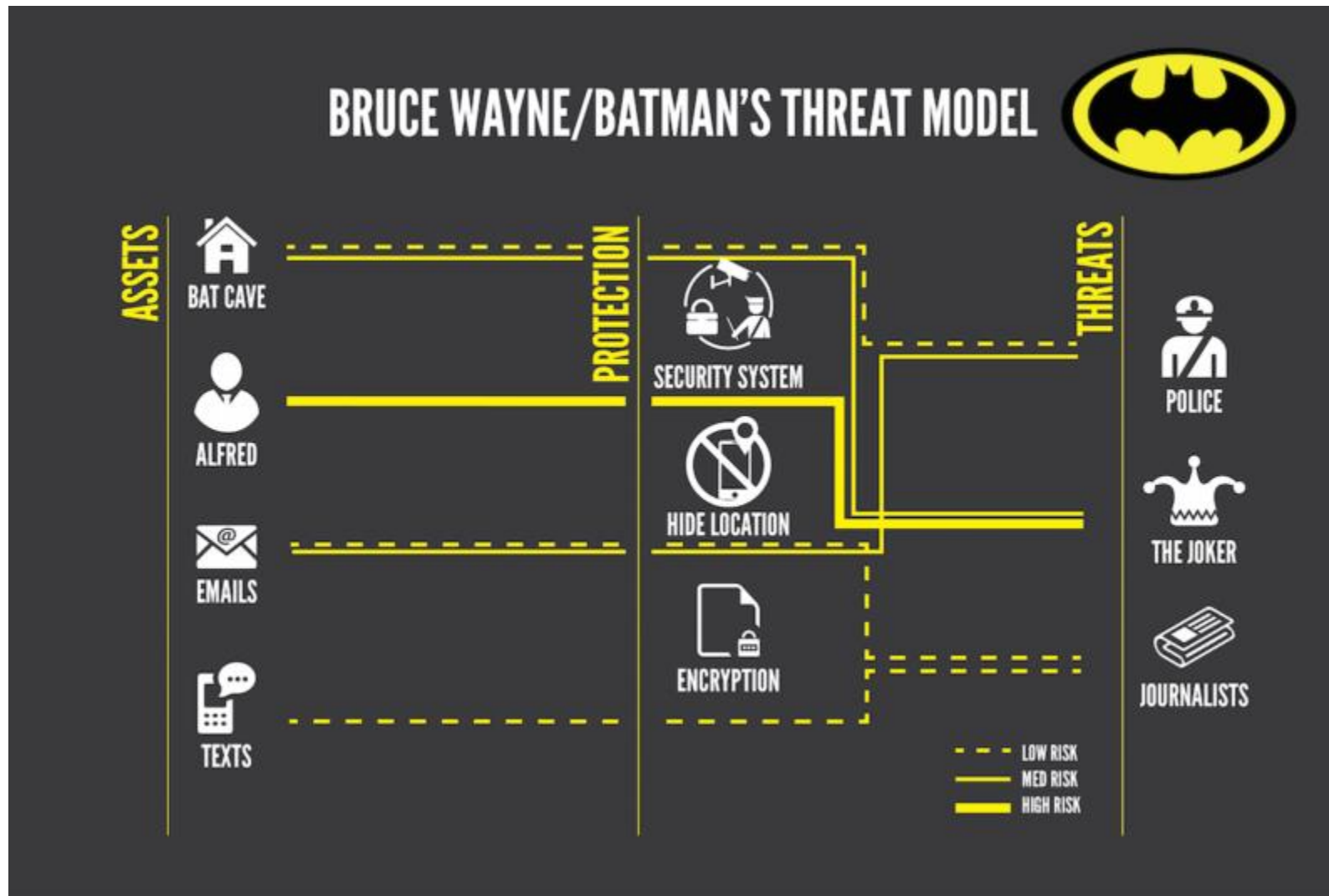
- **Adversary Goals:** What is the adversary attempting to do?
- **Adversary Capabilities:** What resources can the adversary bring to bear to accomplish their goals?

SECURITY MEANS PREVENTING GOALS FROM BEING ACCOMPLISHED,
DESPITE CAPABILITIES

- **Defender Capabilities:** What resources **MUST** be brought to bear to defeat the threat model?

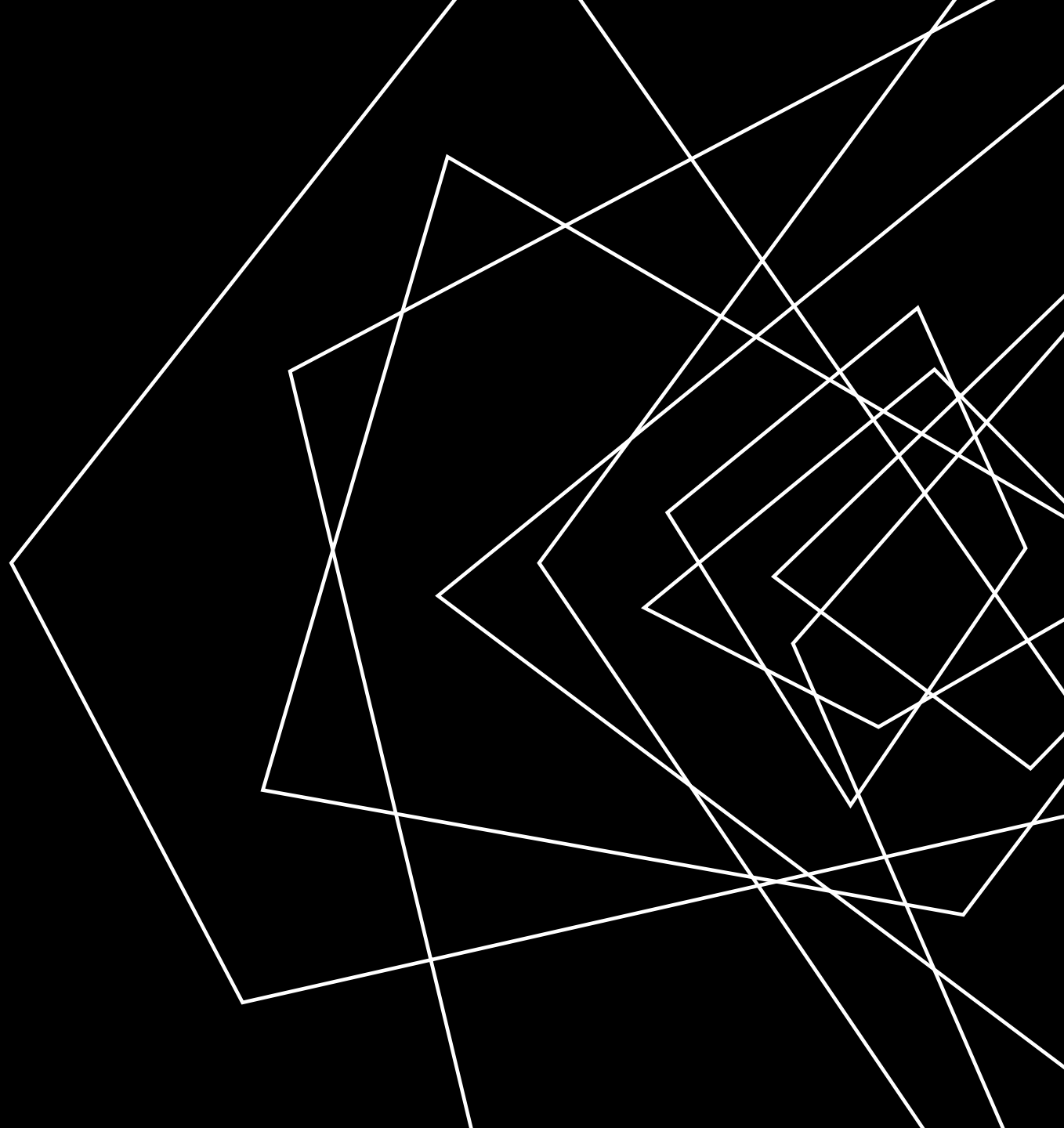
THINKING ABOUT ATTACKS

THREAT MODELS



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THE BASIC IDEA OF SIDE CHANNELS

SIDE CHANNELS

ABSTRACTION IS A KEY PRINCIPLE OF COMPUTER SCIENCE!

As a programmer, you shouldn't need to know underlying details

AS A SECURITY EXPERT, THESE DETAILS MIGHT END UP BEING IMPORTANT!

The way a program accomplishes its tasks are important, especially from a security aspect

- How long does it take for the program to do X ?
- How hot does it make the processor when X happens?
- How much power does it draw when X happens?

SIDE CHANNELS – THE BIG IDEA

SIDE CHANNELS - INSTANCES

COMPUTATION MAY HAVE EFFECTS OUTSIDE OF PROGRAM SEMANTICS

Some operations (internally) take longer based on aspects of the data

TEMPEST

SIDE CHANNELS – HISTORY

ELECTROMAGNETIC LEAKAGE OF KEYS

- **WWII:** Bell Telephone discovers electromagnetic leakage in one-time pad teleprinters, detectable at 100-ft radius
- **1951:** CIA rediscovers leakage, detectable at 200-ft radius
- **1964:** TEMPEST shielding protocol established



TEMPEST

SIDE CHANNELS – HISTORY

ELECTROMAGNETIC LEAKAGE OF KEYS

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- **1964:** TEMPEST shielding protocol established



VAN ECK PHREAKING

SIDE CHANNELS – HISTORY

ELECTROMAGNETIC LEAKAGE OF MONITORS

- Pick up the monitor's electromagnetic emanations that differ depending on how the screen lights up
- Originally determined for CRT (1985), also discovered for LCD monitors (2004)

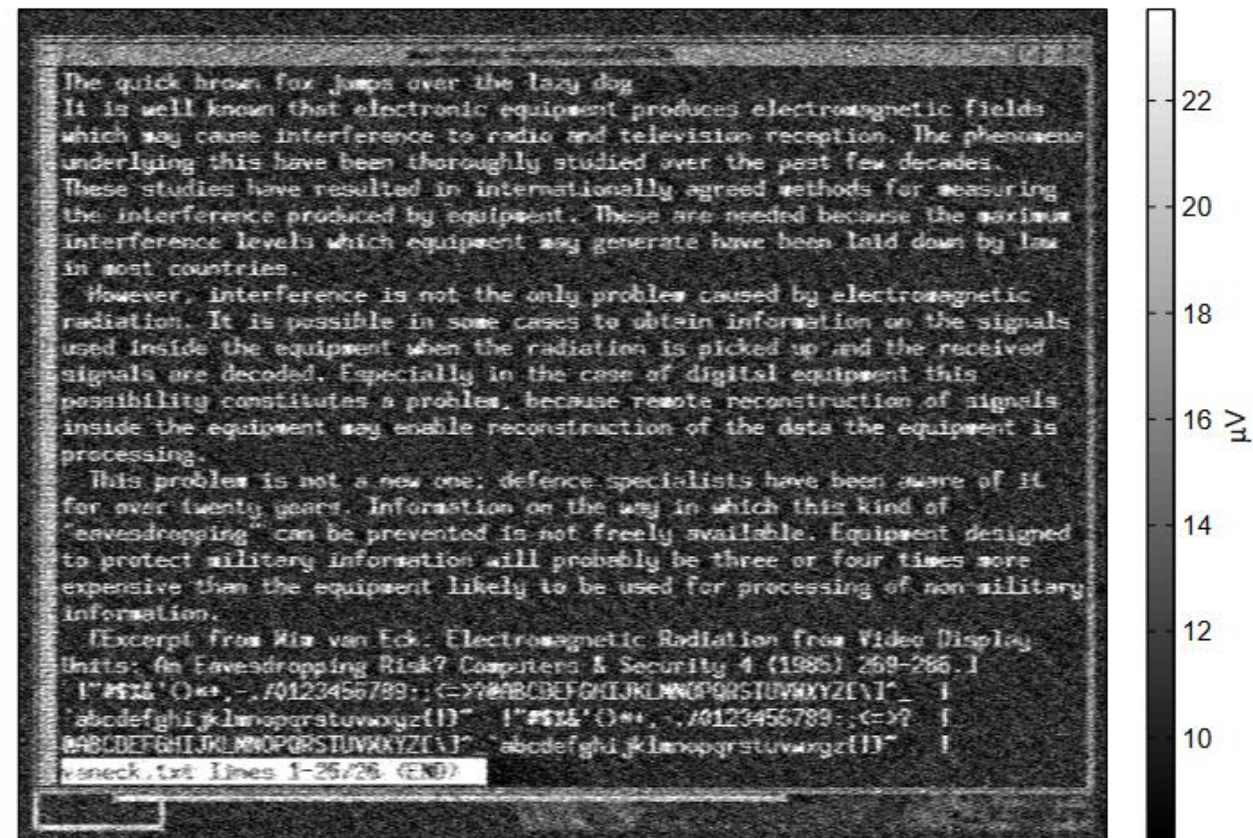


Fig. 3. Text signal received from a 440CDX laptop at 10 m distance through two intermediate offices (3 plasterboard walls).

SIDE CHANNELS – PARTIAL CREDIT

SIDE CHANNELS - INSTANCES

EVEN “HINTS” ABOUT SECRET DATA CAN BE PROBLEMATIC

Assume you’re trying to guess a password

- knowing even 1 character massively reduces the search space
- knowing the length of the password reduces the search space



Partial
Credit

COVERT CHANNELS

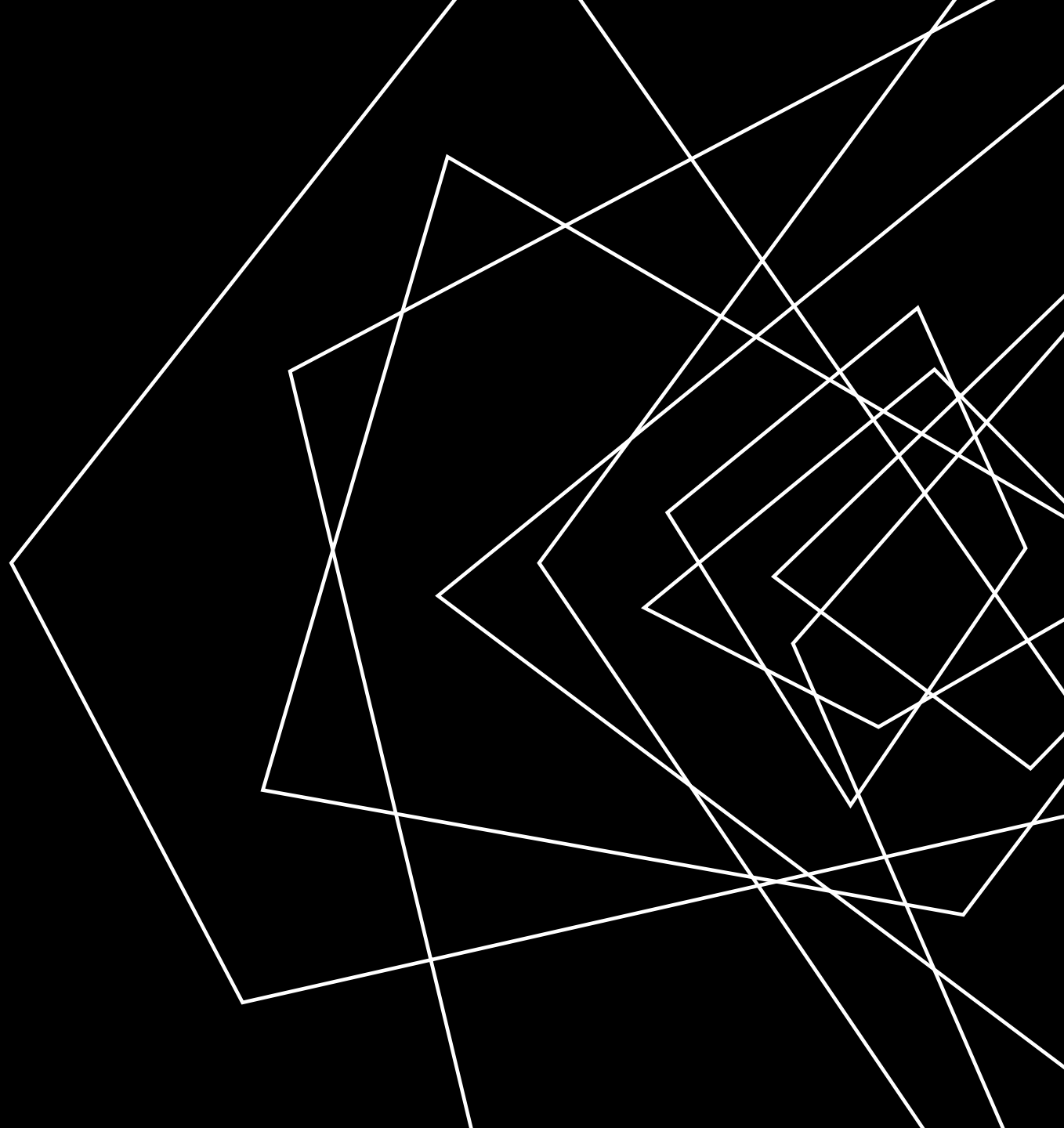
SIDE CHANNELS

SOMETIMES A PROGRAM WANTS TO LEAK DATA

Exfiltration !

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TIMING SIDE CHANNELS

SIDE CHANNELS - INSTANCES

SOME COMPUTATIONS TAKE LONGER THAN OTHERS

Some operations (internally) take longer based on aspects of the data

```
bool checkPW(const char * given){
    const char * expected = "12345";
    int len = min(5, strlen(given));
    for (int i = 0; i < len, i++){
        if (given[i] != expected[i]){
            return false;
        }
    }
    return true;
}
```

TIMING SIDE CHANNELS

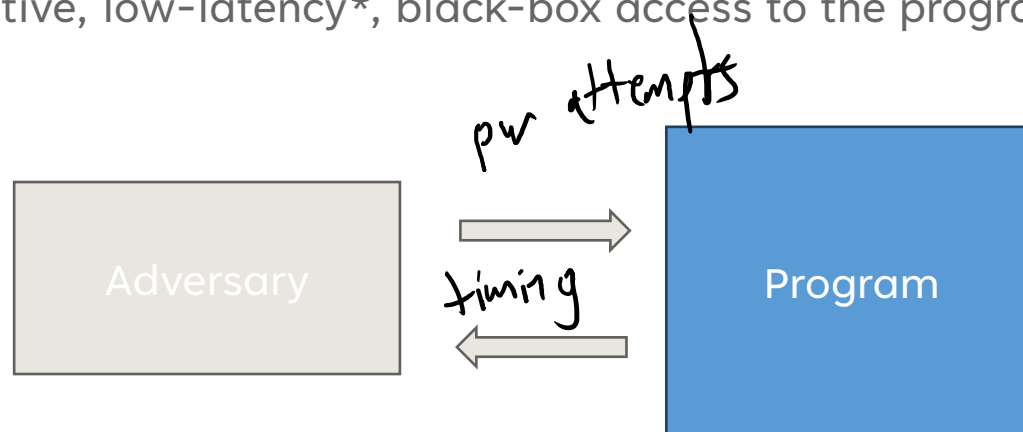
SIDE CHANNELS - INSTANCES

SOME COMPUTATIONS TAKE LONGER THAN OTHERS

Some operations (internally) take longer based on aspects of the data

THREAT MODEL

Interactive, low-latency*, black-box access to the program, precise timer



*: May be overcome with more samples

TIMING SIDE CHANNELS - FIX

SIDE CHANNELS - INSTANCES

```

bool checkPW(const char * given){
    const char * expected = "12345";
    int len = min(5, strlen(given));
    for (int i = 0; i < len, i++){
        if (given[i] != expected[i]){
            return false;
        }
    }
    return true;
}

```

bool ok = true;

```

bool checkPW(const char * given){
    const char * expected = "12345";
    int len = min(5, strlen(given));
    for (int i = 0; i < len, i++){
        if (given[i] != expected[i]){
            return false; ok = false;
        }
    }
    ok
    return true;
}

```

TIMING SIDE CHANNELS - FIX

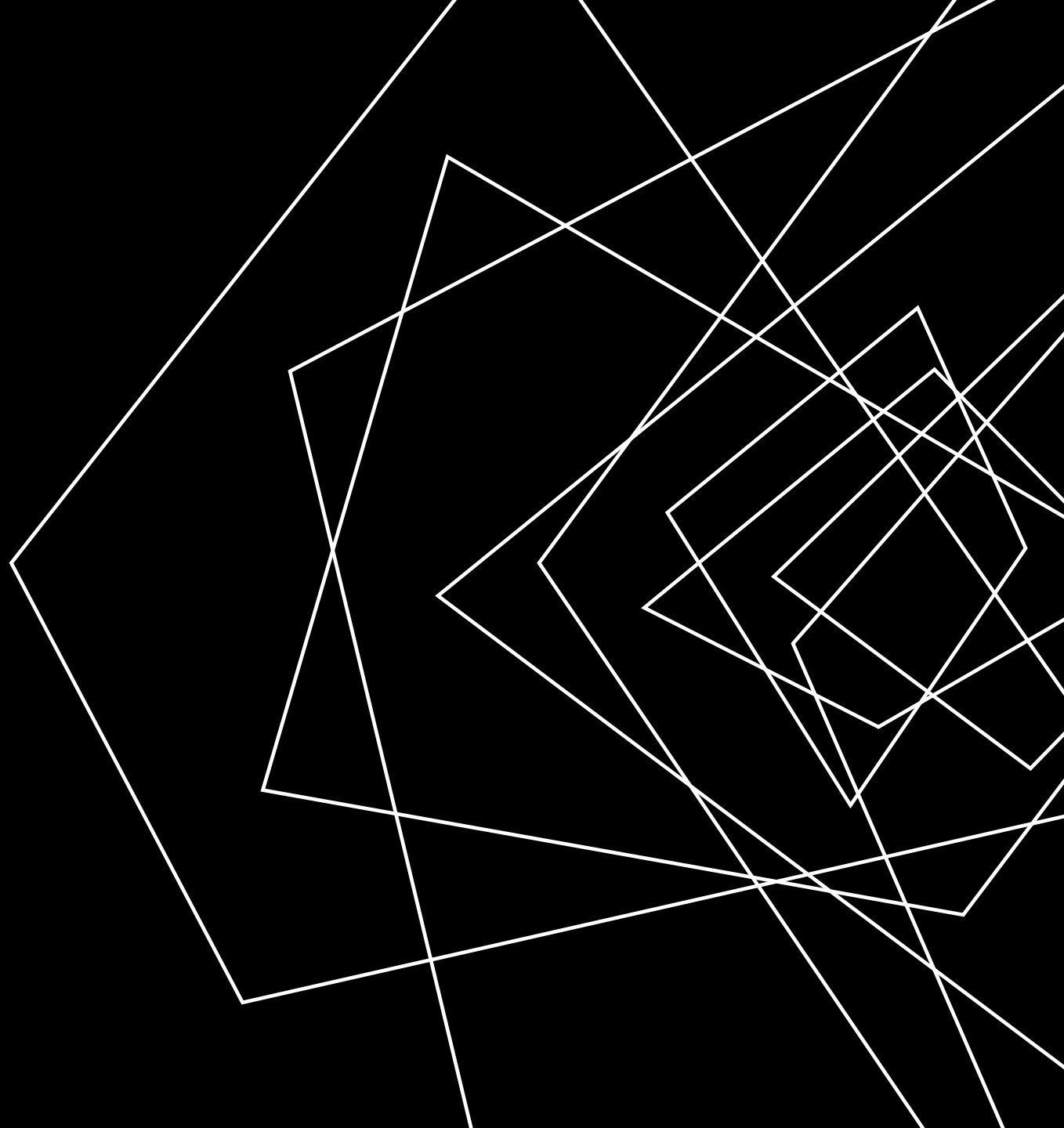
SIDE CHANNELS - INSTANCES

LIMITATIONS OF UNIFORM EXECUTION

- Necessarily slow down your computation to the worst case
- May require some pretty precise understanding of timing
- May not always be obvious what the worst-case even is

LECTURE OUTLINE

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- Instances
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TIMING SIDE CHANNELS - FIX

SIDE CHANNELS - INSTANCES

CAN WE FIX THIS ISSUE WITH OUR DATAFLOW APPROACH?

- Instruction transformers: how much time that instruction takes
- Block composition: the sum total of instruction times
- Merge operation: some sort of check that all paths are of comparable time?

WRAP-UP

